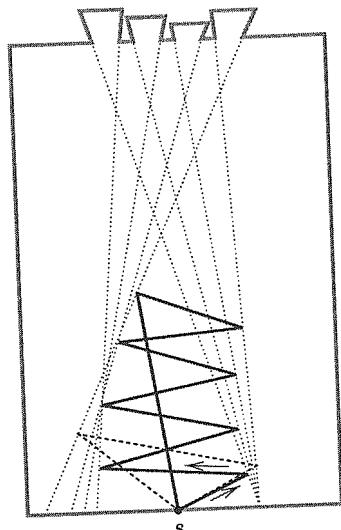
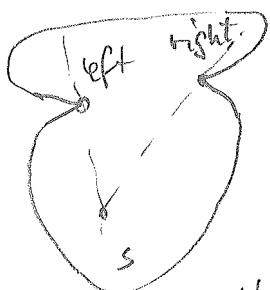


Exploring cuts in clockwise
order doesn't work in
non-rectilinear polygons.



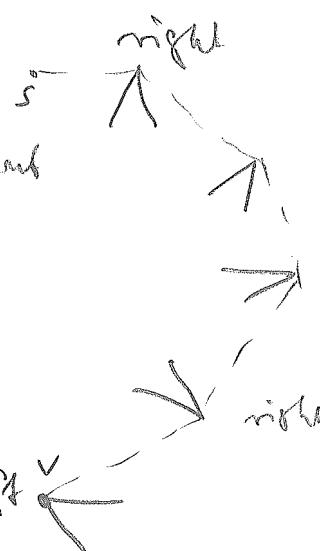
Idea Explore left and right
reflex vertices separately.



Problem

How to know about
left vertex v ?

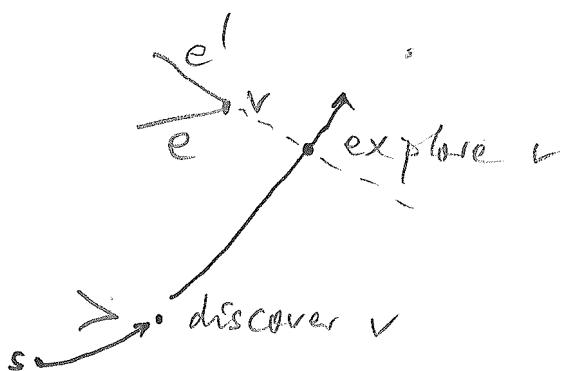
Idea Partition left/right reflex
vertices into groups
start with right vertices



notations:

reflex vertex v discovered:
(part of) edge e has been seen

... explored: (part of) edge e
has been seen, too



CP: robot's current position

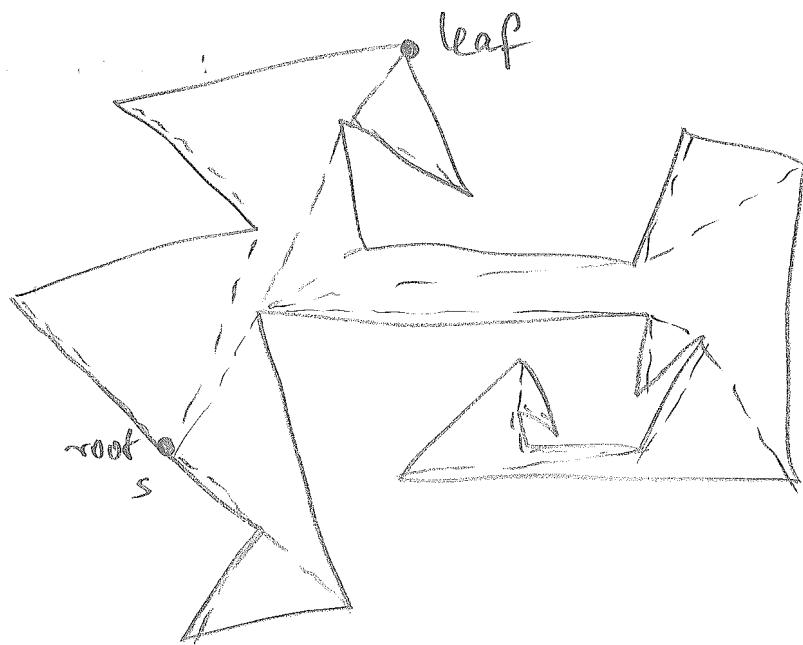
TargetList: list of discovered, but not yet explored,
right vertices, sorted in clockwise order on γ
(initially: all right vertices discoverable, but
not explorable, from s)

BasePoint: point where search algorithm starts
(initially = CP)

another useful structure:

shortest path tree SPT(s)

- tree made of all shortest paths from s to the vertices of P



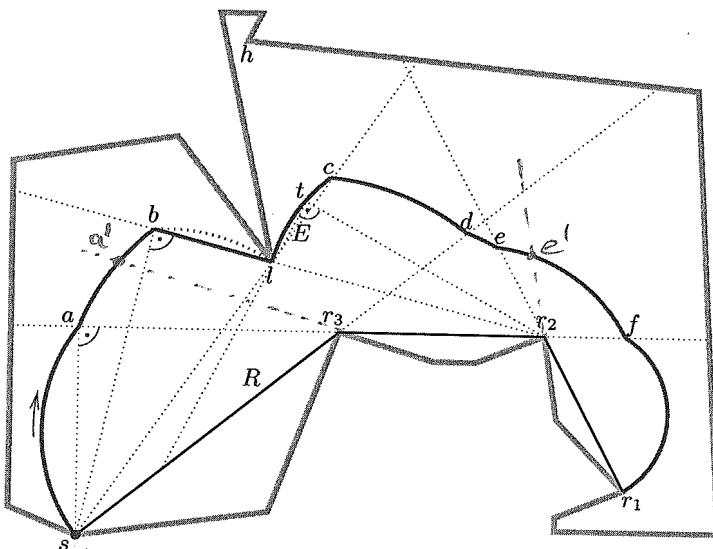
internal nodes

- reflex vertices

```

procedure ExploreRightVertex ( inout TargetList, inout ToDoList );
  BasePoint := CP;
  Target := First (TargetList);
  if Target not visible then
    walk on shortest path from BasePoint to Target
    until Target becomes visible;
  Back := last vertex before CP on shortest path from BasePoint to CP;
  walk clockwise along circ (Back, Target)
  while maintaining TargetList and ToDoList
    whenever First (TargetList) changes let Target := First (TargetList);
    whenever Back becomes invisible update Back;
  exceptions for walking along the circle:
    if the boundary of P blocks the walk on the current circle then
      walk clockwise along the boundary
      until the circular walk is again possible;
    if Target is becoming invisible then
      walk towards Target
      until the blocking vertex is reached;
  until Target is fully explored;
end ExploreRightVertex;

```



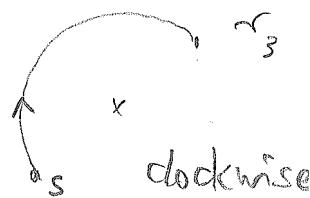
Discussion of Example:

Initially, BasePoint = CP = s ,
 TargetList = $\{r_3\}$, ToDoList = \emptyset

Target = r_3 is visible

Back = s

robot starts following circ (s, r_3)

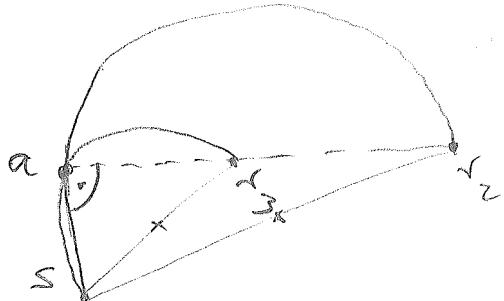


at point a, right reflex vertex γ_2 is discovered
comes before γ_3 in clockwise order

$$\Rightarrow \text{Target} = \gamma_2$$

robot now follows $\text{arc}(s, \gamma_2)$

observe: $\text{arc}(s, \gamma_2)$ does pass through a
(because of 'Thales' theorem)



at point a' , γ_3 gets fully discovered, removed from Target list
at point b, Target γ_2 would become invisible

\Rightarrow robot walks towards Target γ_2 until blocking vertex l is reached

robot now follows $\text{arc}(l, \gamma_2)$ because
 $\text{Back} = l$ = last vertex of shortest path
from BasePoint s to CP

at point c, BasePoint s becomes visible again

\Rightarrow robot now follows $\text{arc}(s, \gamma_2)$ because
 $\text{Back} = s$

at point , BasePoint s becomes invisible

robot follows $\text{arc}(\gamma_3, \gamma_2)$ because
 $\text{Back} = \gamma_3$ = last vertex on shortest path
from BasePoint s to CP

at e, right reflex vertex r_1 is discovered
comes before r_2 in clockwise order

$\Rightarrow \text{Target} = r_1$
robot follows circ(r_2, r_1)

at e' , r_2 becomes fully explored and removed
from Target List

at f, r_3 becomes invisible

robot now follows circ(r_2, r_1) because

Back = r_2 = last vertex on shortest path
from BasePoint s to CP

at r_1 , Target r_1 becomes fully explored.

Observation Robot reaches cut of Target at
the same point as shortest path from BasePoint

Remarks

(i) right reflex vertices r_2, r_1 were added to
Target List because shortest paths from s
to r_2, r_1 did not contain left turns

(ii) right reflex vertex h was not added to
Target List, because shortest path s - h
makes left turn at l.

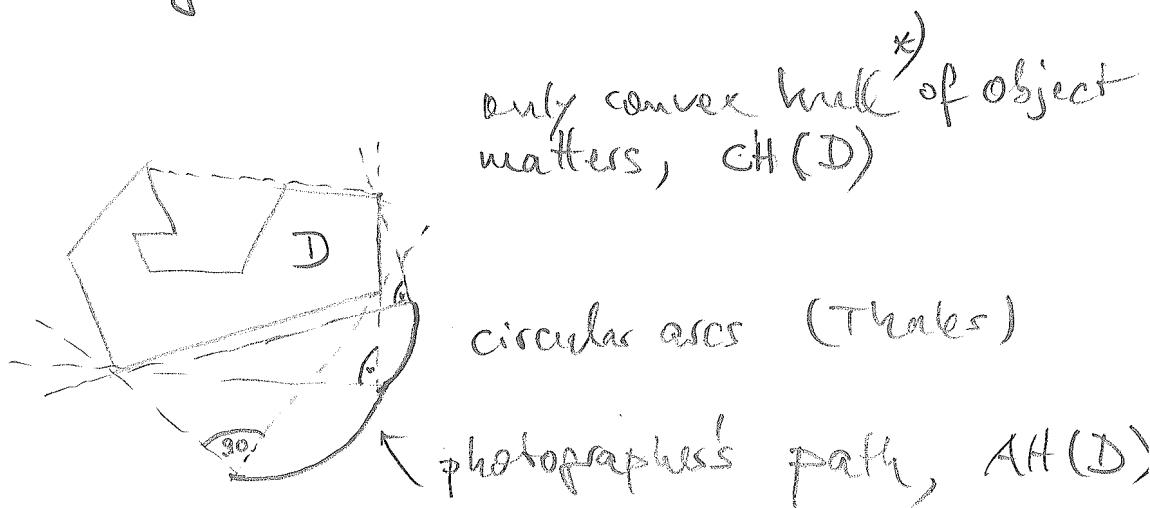
reflex vertex h will be listed later on (To Do List)

Question: How long is path generated by
ExploreRightVertex
as compared to shortest path from s to r_1
(in example)?

need structural property for analysis

Photographer's path \triangleq angle hull

Suppose photographer moves around object to find best view point; at all times object fits exactly into 90° lens angle



Lemma 1 in the free plane,

$$|AH(D)| \leq \frac{\pi}{2} \cdot |(CH(D))| \leq \frac{\pi}{2} |D|,$$

where $|.|$ denotes the perimeter.

(without proof)

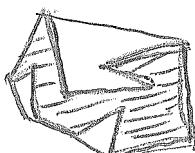
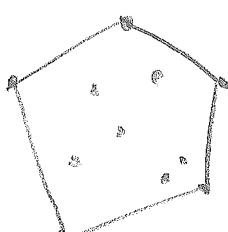
Example:



$$\text{perimeter}(D) = 2\pi r$$
$$\text{perimeter}(AH(D)) = \pi d$$

We are interested in a constrained case where object D is located in simple polygon P

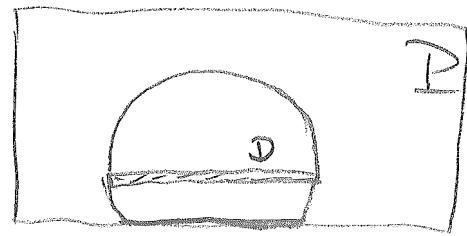
* convex hull $CH(D) := \bigcap_{\substack{C \\ D \subseteq C \subseteq \mathbb{R}^2, \\ C \text{ convex}}} C$



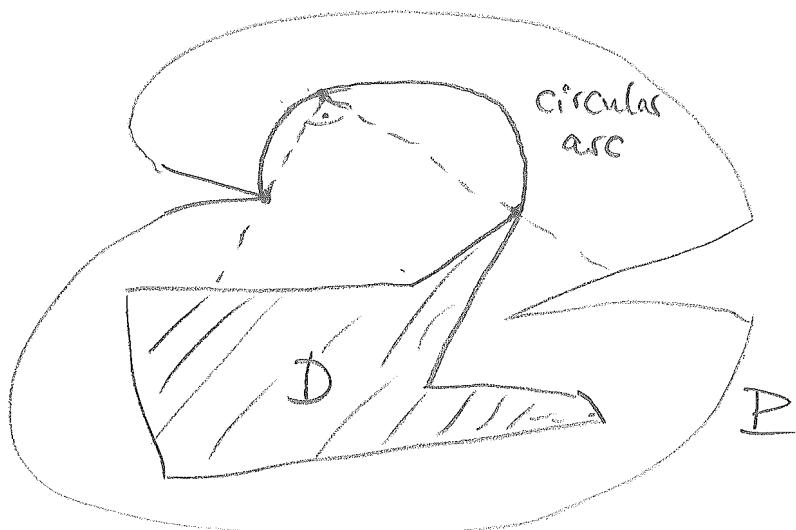
examples

Here, photograph's path may

- follow a polygon's edge ("too close")



- be constrained by reflex polygon vertex touching 90° wedge from the outside



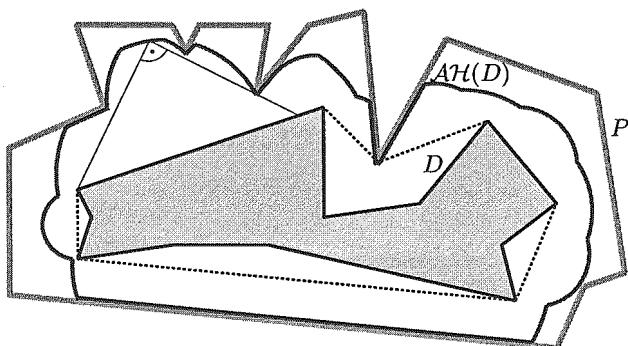
more precisely,

Def: Let D be a connected compact set in the interior of a simple polygon P . Then,

$$AH_{\frac{\pi}{2}}(D) := \{ z \in P \mid z \text{ can see two points of } D \text{ at a } 90^\circ \text{ angle} \}$$

is called the angle hull of D in P .

(Its boundary : photograph's path.)



Similar to previous case: boundary of angle hull depends only on relative convex hull $RCH_P(D)$ of D in P (rubber band encircling D in P)

Theorem A $|AHT(D)| \leq 2 \cdot |RCH(D)| \leq 2 \cdot |D|$,
and this bound can be attained.

Proof later

Consequence:

Lemma 2 Path generated by ExploreRightVertex from BasePoint to cut of Target explored is at most twice as long as shortest path.

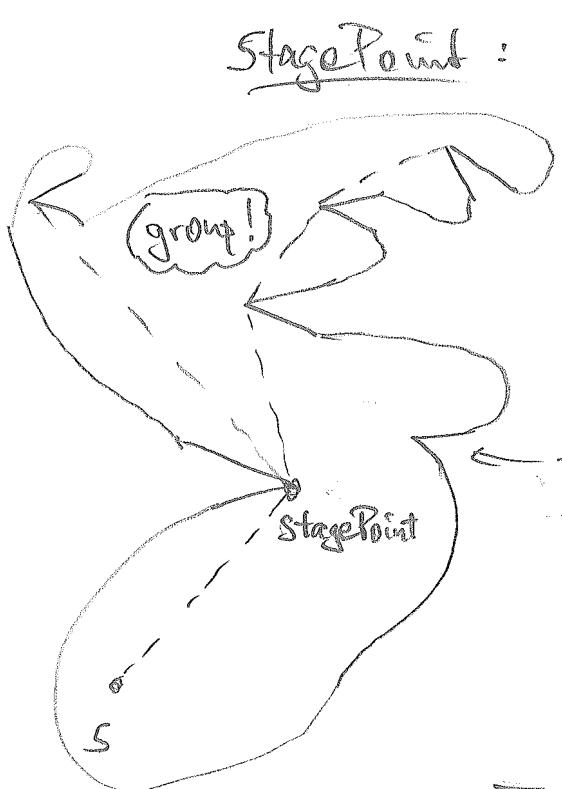
Proof Path generated is part of $\mathcal{D}AHT$ (shortest path), (except for straight segments leading to blocking vertices, which are replaced by longer circular segments in AHT).

Symmetric to ExploreRightVertex: ExploreLeftVertex

Next: exploring a group of right vertices

```
procedure ExploreRightGroup (in TargetList, out ToDoList);
  StagePoint := CP;
  ToDoList := empty list;
  while TargetList is not empty do
    ExploreRightVertex (TargetList, ToDoList);
    (* CP is now on the cut, C, of the last target. *)
    walk to the point on C that is closest to StagePoint
      while maintaining TargetList and ToDoList;
    walk on the shortest path back to StagePoint;
end ExploreRightGroup;
```

In ExploreRightGroup :



initially, $s = s$
in general, a left reflex vertex
of the shortest path tree of s ,
also visited by W_{opt}

explored in earlier stage

On calling ExploreRightGroup :

- To Do List = \emptyset
- TargetList contains sorted list of unexplored right vertices, whose shortest paths from StagePoint make only right turns among them all right children of StagePoint

During execution of ExploreRightVertex (or walkc) :

- all newly discovered right vertices are added to TargetList whose shortest paths from StagePoint contain only right turns
- all explored right vertices are added to To Do List that have a left child
(children of explored vertices are known!)

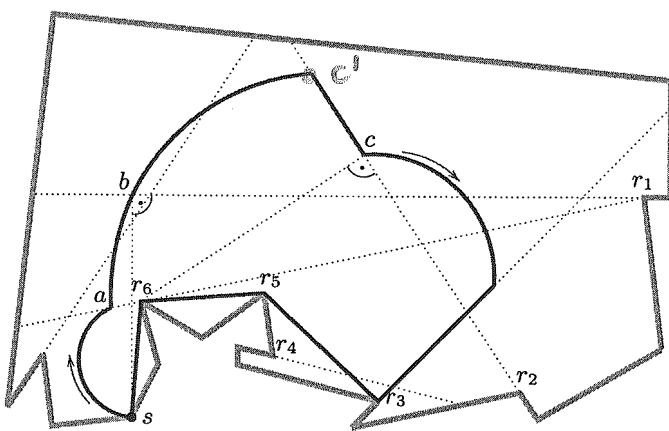
On returning from ExploreRightVertex :

- robot adjusts its position on the cut of last Target (will become useful later)
adjusted points because StagePoint

On termination of ExploreRightGroup:

- all right vertices initially in TargetList, and all their purely right descendants in the shortest path tree of s have been explored.

Example:



during exploration of r_6 , vertex r_1 is discovered, at a
at b , exploration of Target r_1 is completed.

ExploreRightVertex returns,
no adjustment necessary because b is closest point
to StagePoint s on cut of r_1 .

Target list : r_6, r_1 removed, r_2, r_5 added

next call to ExploreRightVertex explores r_2 ,
finishes at c'

robot walks along cut to c = closest point to
StagePoint s

next call to ExploreRightVertex explores r_3 ,

robot walks along cut to c' = closest point to s
while doing so, r_4 gets explored

robot returns on shortest path to StagePoint, s .