

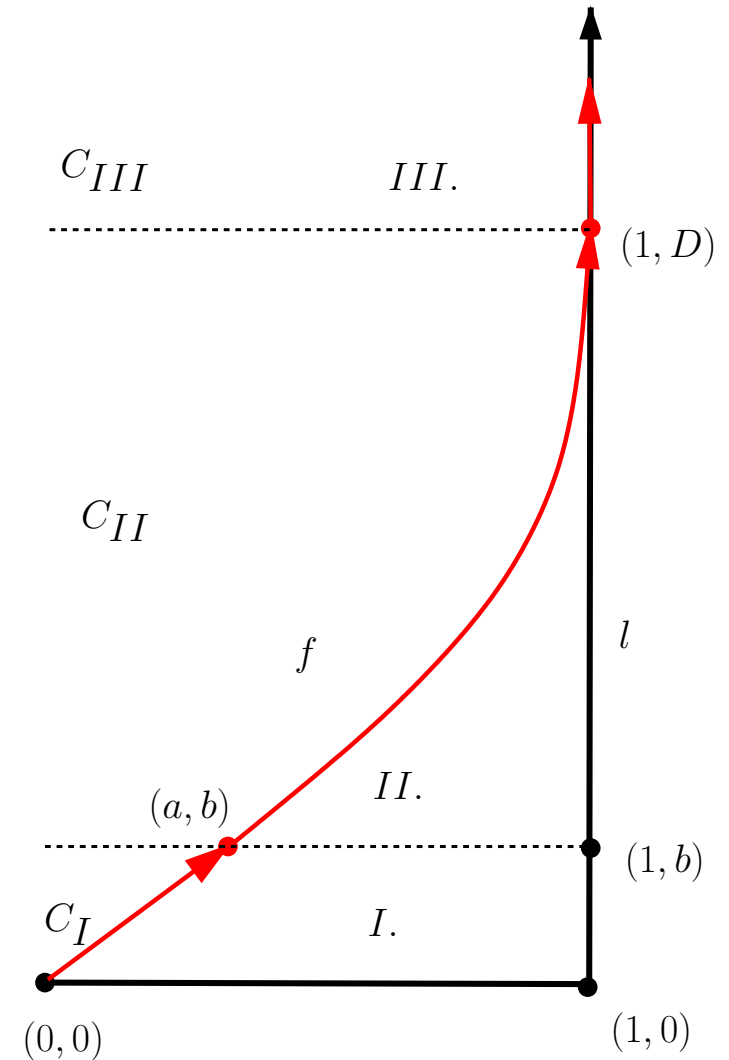
Online Motion Planning MA-INF 1314

General rays!

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Rep.: Window Shopper Strategy: Three parts

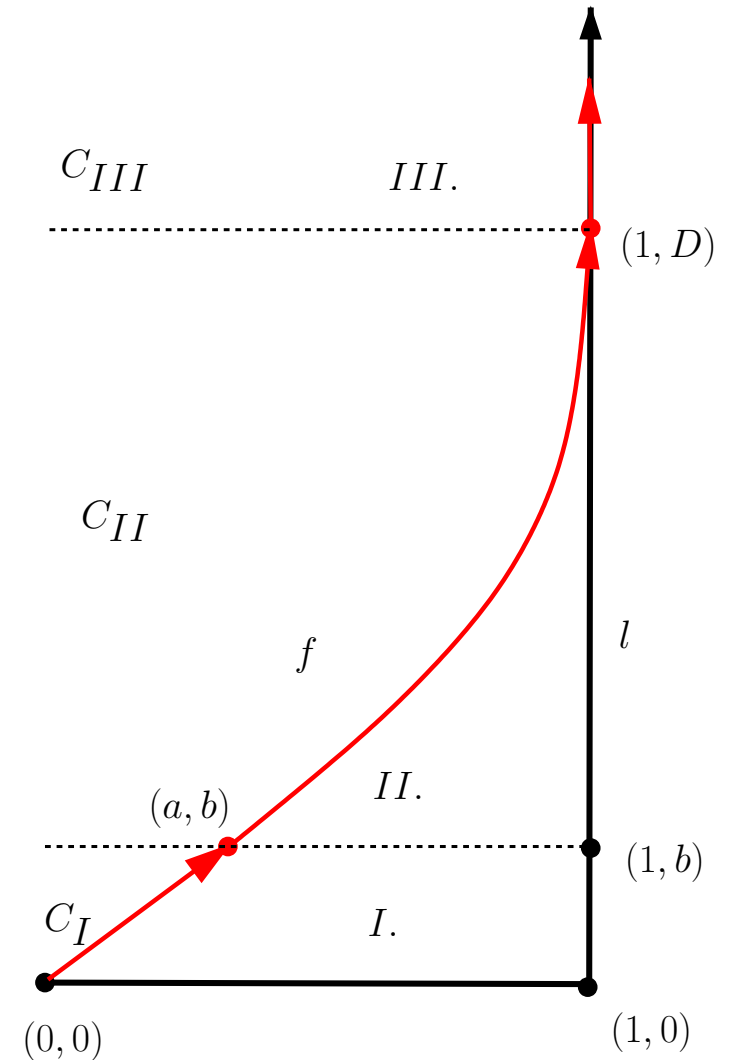
- A line segment from $(0, 0)$ to (a, b) with **increasing** ratio for s between $(1, 0)$ and $(1, b)$
- A curve f from (a, b) to some point $(1, D)$ on l which has **the same** ratio for s between $(1, b)$ and $(1, D)$
- A ray along the *window* starting at $(1, D)$ with **decreasing** ratio for s beyond $(1, D)$ to infinity
- Worst-case ratio is attained for all s between $(1, b)$ and $(1, D)$
- Optimality by construction, CONVEX!



Rep.: Design of the strategy

- Condition I): $a = \frac{1-b^2}{2}, \sqrt{1+b^2}$ ■
- Condition II):

$$f'(x) = 2\sqrt{1+b^2} \frac{\sqrt{1+f(x)^2} f(x)}{1-b^2 f(x)^2},$$
 Point $((1-b^2)/2, b)$ ■
- Solve: $y' = 1 \cdot 2\sqrt{1+b^2} \frac{\sqrt{1+y^2} y}{1-b^2 y^2}$ with
 $((1-b^2)/2, b)$ ■
- First order diff. eq. $y' = h(x)g(y)$ ■
- Solution: $\int_l^y \frac{dt}{g(t)} = \int_k^x h(z) dz:$
 Here $x = f^{-1}(y)$ ■



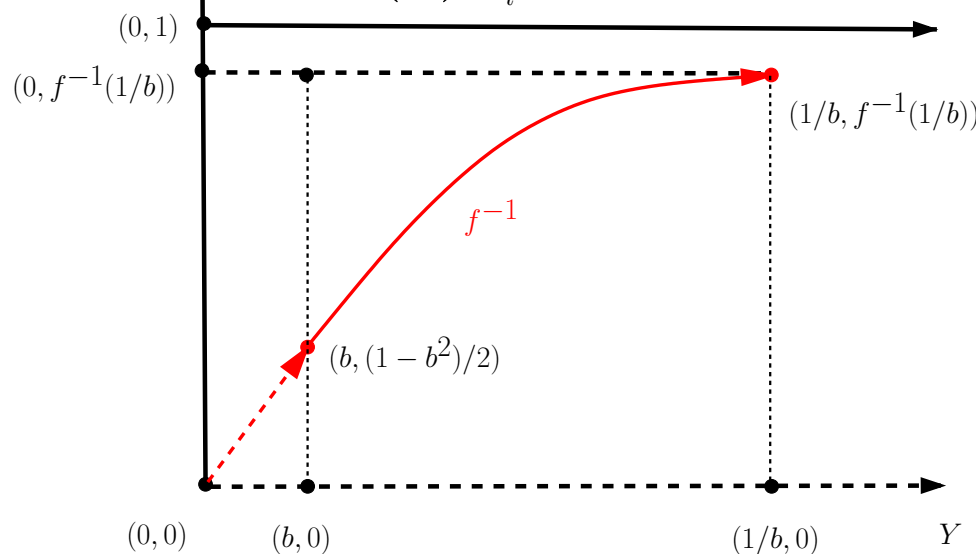
Rep.: Consider inverse function $x = f^{-1}(y)$

- $x = \frac{b^2\sqrt{1+y^2} + \operatorname{arctanh}\left(\frac{1}{\sqrt{1+y^2}}\right) - \operatorname{arctanh}\left(\frac{1}{\sqrt{1+b^2}}\right) - \sqrt{1+b^2}}{2\sqrt{1+b^2}}$

- $x' = \frac{1}{g(y)} = -\frac{(b^2y^2-1)}{2\sqrt{1+y^2}y\sqrt{1+b^2}} \geq 0$ for $y \in [b, 1/b]$

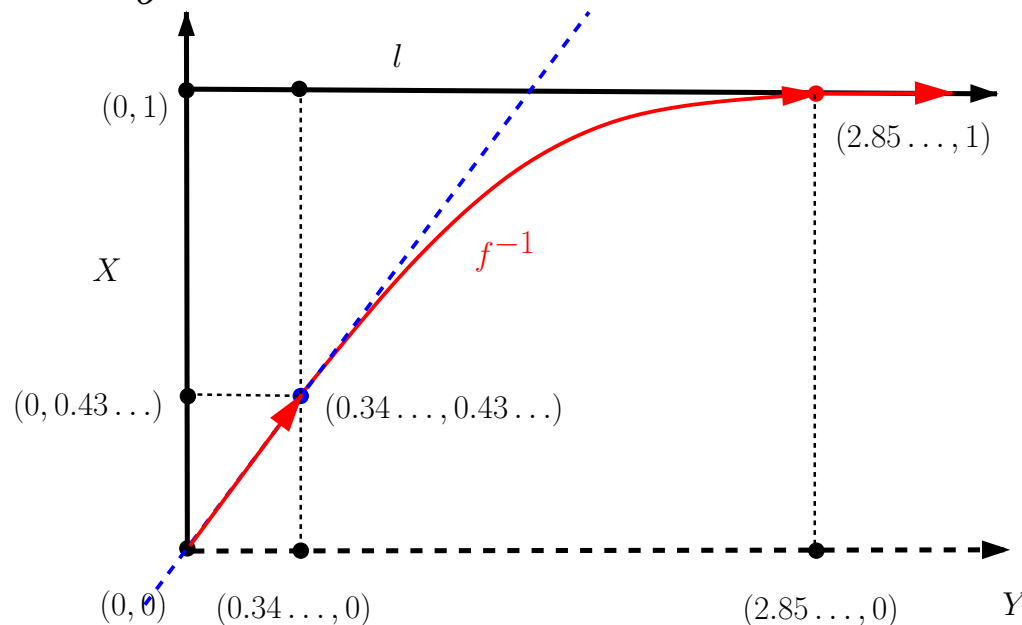
- $x'' = -\frac{(b^2y^2+2y^2+1)}{2(1+y^2)^{3/2}\sqrt{1+b^2}y^2} \leq 0$ for $y \geq 0$

- $x = f^{-1}(y)$ concave, $y = f(x)$ convex, Max. at $y = 1/b$



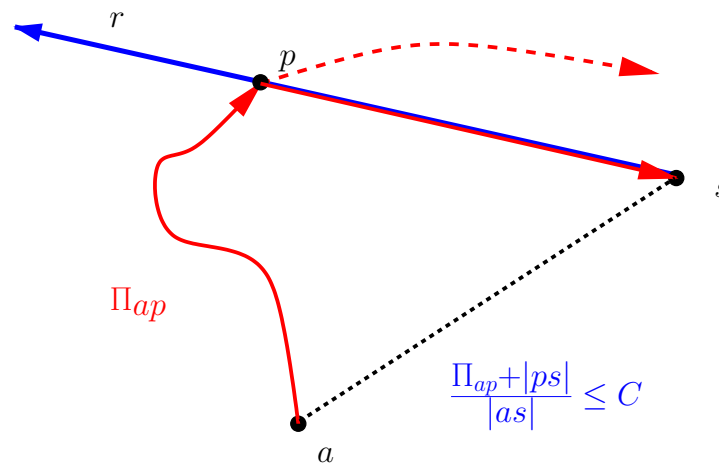
Rep.: Theorem Optimality of f

- Solve $f^{-1}(1/b) = 1$: $b = 0.3497\dots$, $D = 1/b = 2.859\dots$,
- $a = 0.43\dots$, ■ worst-case ratio $C = \sqrt{1 + b^2} = 1.05948\dots$ ■
- f convex from (a, b) to $(1, D)$, line segment convex■
- Prolongation of line segment is tangent of f^{-1} at (b, a) ■
- Insert: $f^{-1}'(b) = \frac{a}{b}$! Convex!■



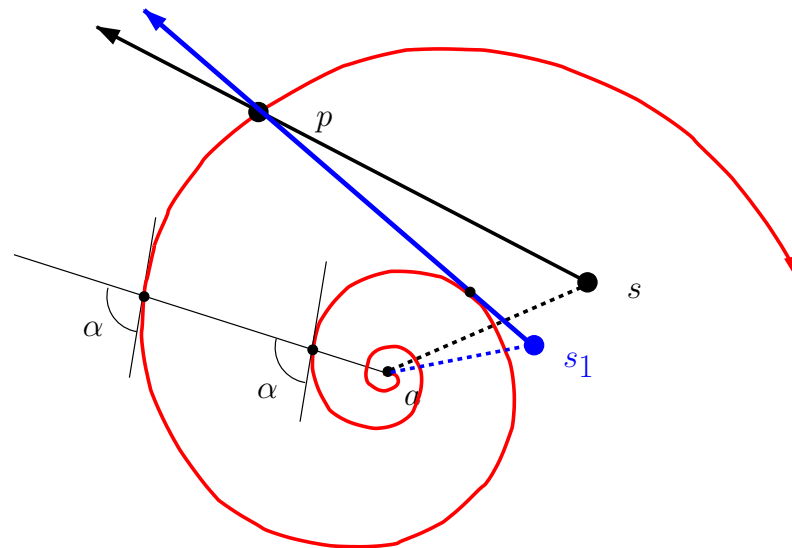
Rep.: Rays in general

- Rays are somewhere in the plane
- ● Searchpath Π
- Upper bound: $C = 22.531 \dots$
- Lower bound: $C \geq 2\pi e = 17.079 \dots$



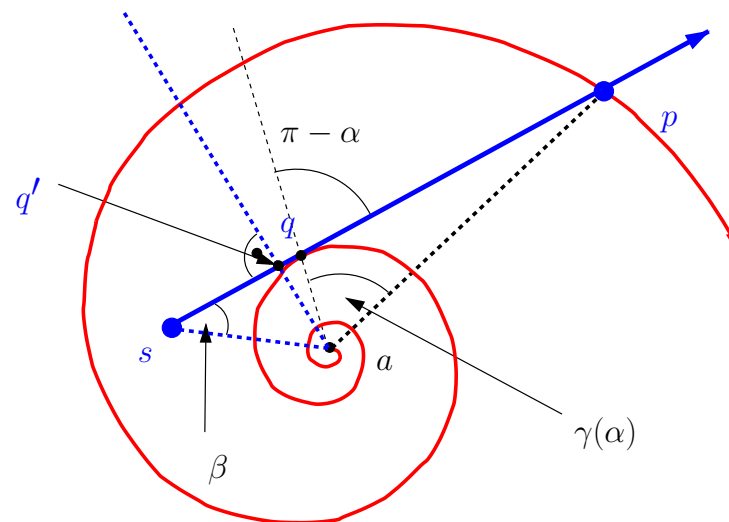
Rep.: Strategy, Spiral search

- Logarithmic spiral
- Polar coordinates $(\varphi, d \cdot E^{\phi \cot(\alpha)})$, $d > 0, -\infty < \varphi < \infty$
- $\alpha = \pi/2$ Kreis!
- Hits the ray, moves to s
- Worst-case ratio: Ray is a tangent **Lemma**



Rep.: Optimizing the spiral

- Strategy $d \cdot E^{\phi \cot(\alpha)}$, property $|\text{SP}_a^p| = \frac{|ap|}{|\cos(\alpha)|} = \frac{dE^{\theta_p \cot \alpha}}{|\cos(\alpha)|}$ ■
- Ratio C identical for all tangents: Ratio $C(\alpha)$ ■
- We optimize for perpendicular points q' ■
- Adversary can move s a bit to the left (chooses β) ■
- Law of sine: Ratio $C(\beta, \alpha) = C(\alpha) \sin(\beta) + \cos(\beta)$ ■

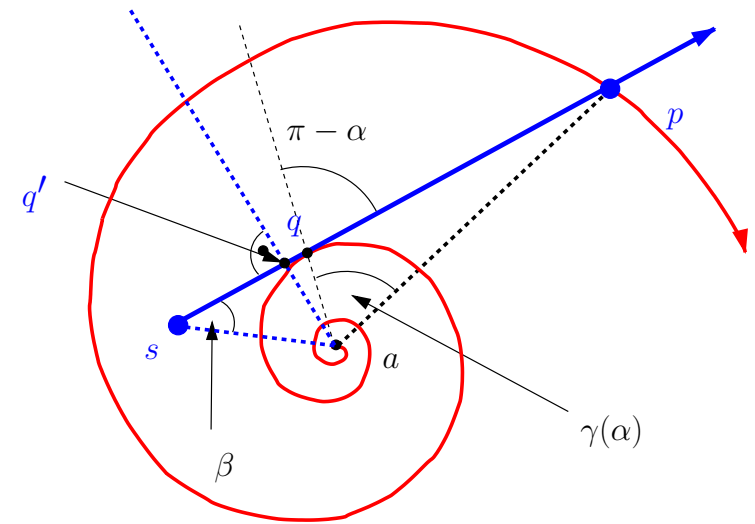


Rep.: Minimize worst-case ratio

- $dE^{\phi \cot(\alpha)}$: Determine α , $d = 1$ ■
- Assume: Fixed α for q' ■
- $|as| \sin(\beta) = |aq'|$,
 $|sq'| = |as| \cos(\beta) = \frac{|aq'| \cos(\beta)}{\sin(\beta)}$ ■
- Ratio $C_{q'} = \frac{|SP_a^p| + |pq'|}{|aq'|}$
 maximized by

$$\frac{C_{q'} \cdot |aq'| + |aq'| \cos(\beta) / \sin(\beta)}{|aq'| / \sin(\beta)} =$$

$$C_{q'} \sin(\beta) + \cos(\beta)$$
 ■
- Minimize over α for q' ■
- Finally adversary choose β , fixes s ! ■



Minimize worst-case for q'

- $p = (\phi, E^{\phi \cot(\alpha)})$: Determine $\alpha < \pi/2$

- $|SP_a^p| = \frac{|ap|}{\cos(\alpha)} = \frac{E^{\cot \alpha (2\pi + \gamma + \theta_q)}}{\cos(\alpha)}$

- $|pq| \sin(\alpha) = |ap| \sin(\gamma)$ and

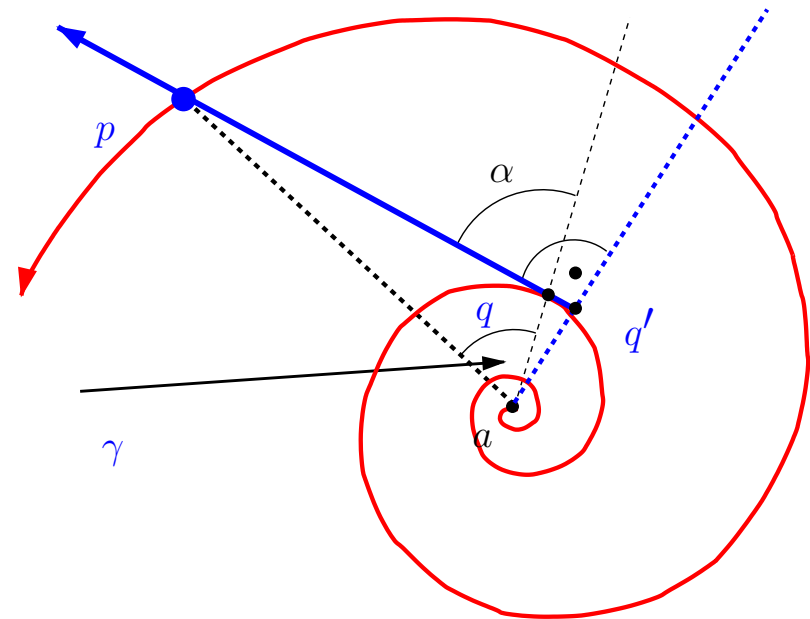
- $|pq| = \frac{E^{\cot \alpha (2\pi + \gamma + \theta_q)} \sin(\gamma)}{\sin \alpha}$

- $|qq'| = |aq| \cos(\alpha)$

- $|pq'| = \frac{E^{\cot \alpha (2\pi + \gamma + \theta_q)} \sin(\gamma)}{\sin \alpha} + E^{\cot \alpha (\theta_q)} \cos(\alpha)$

- $|aq'| = |aq| \sin(\alpha) = E^{\cot \alpha (\theta_q)} \sin(\alpha)$

- γ depends only on α : Now $\gamma(\alpha)$



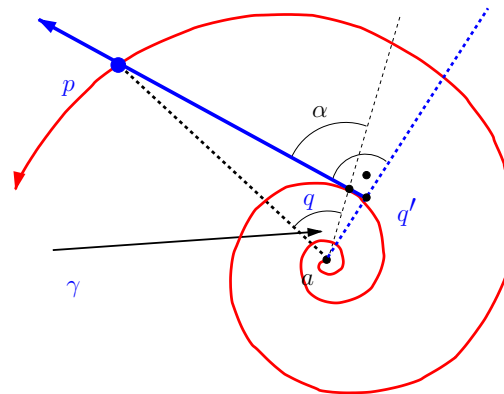
Minimize worst-case for q'

Determine $C_{q'}(\alpha) = \frac{|\Pi_a^p| + |pq'|}{|aq'|}$!

$$\frac{\frac{1}{\cos \alpha} E^{\cot \alpha (\theta_q + 2\pi + \gamma(\alpha))} + \frac{1}{\sin \alpha} E^{\cot \alpha (\theta_q + 2\pi + \gamma(\alpha))} \sin \gamma(\alpha) + E^{\cot \alpha \theta_q} \cos \alpha}{E^{\cot \alpha \theta_q} \sin \alpha} =$$

$$\frac{\frac{1}{\cos \alpha} E^{\cot \alpha (2\pi + \gamma(\alpha))} + \frac{1}{\sin \alpha} E^{\cot \alpha (2\pi + \gamma(\alpha))} \sin \gamma(\alpha) + \cos \alpha}{\sin \alpha} =$$

$$\left(\frac{1}{\sin \alpha \cdot \cos \alpha} + \frac{\sin \gamma(\alpha)}{\sin^2 \alpha} \right) E^{b(2\pi + \gamma(\alpha))} + \cot \alpha$$



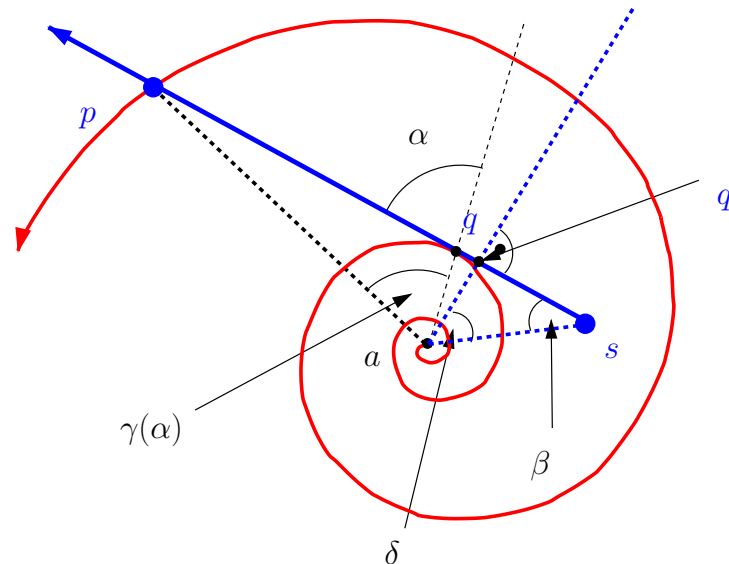
Minimize worst-case for $s = q'$

Determine $\gamma(\alpha)$, Exercise!

Solve Equation: $\frac{\sin \alpha}{\sin(\alpha - \gamma(\alpha))} = E^{\cot \alpha (2\pi + \gamma(\alpha))}$ ■

Then optimize: $f(\alpha) := \left(\frac{1}{\sin \alpha \cdot \cos \alpha} + \frac{\sin \gamma(\alpha)}{\sin^2 \alpha} \right) E^{b(2\pi + \gamma(\alpha))} + \cot \alpha$ ■

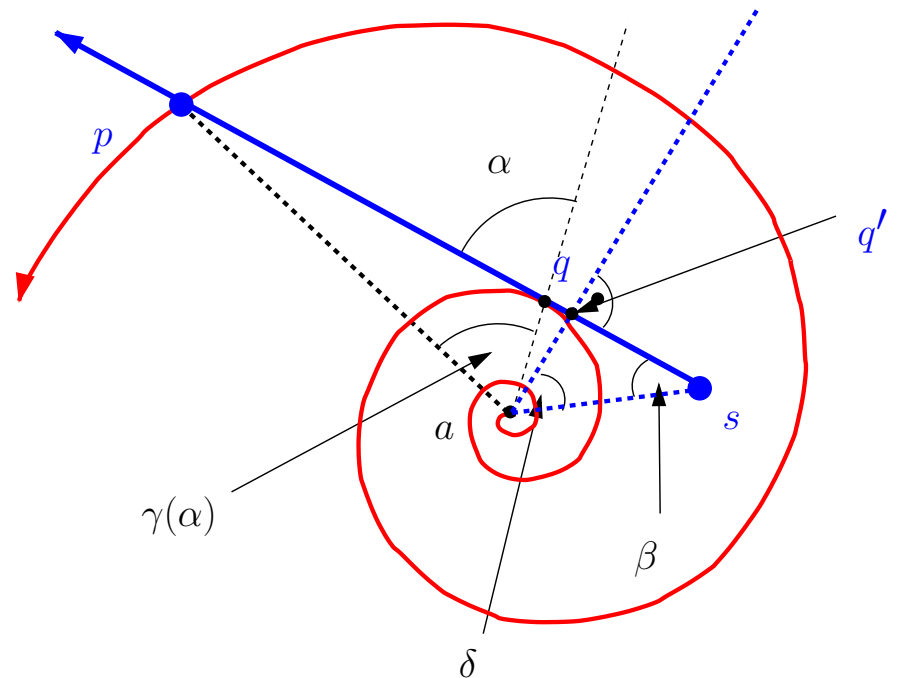
Then minimize: $g(\beta) := f(\alpha_{\min}) \sin \beta + \cos \beta$! ■



Optimizing the spiral: Theorem

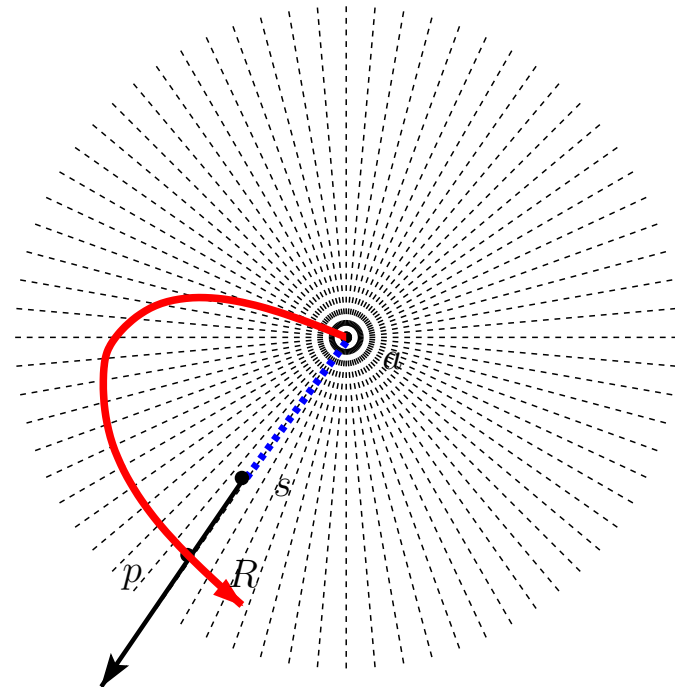
- Ratio: $C(\alpha)$ for $s = q'$ minimal for $\alpha = 1.4575 \dots$ ■
- $C(\alpha) = 22.4908 \dots$ ■
- Adversary choose β for max. $D(\beta, \alpha) = C(\alpha) \sin(\beta) + \cos(\beta)$ ■
- For $\alpha = 1.4575 \dots$ choose $\delta = 0.044433 \dots$ ■

$$D(\beta, C(\alpha)) = 22.51306056 \dots$$



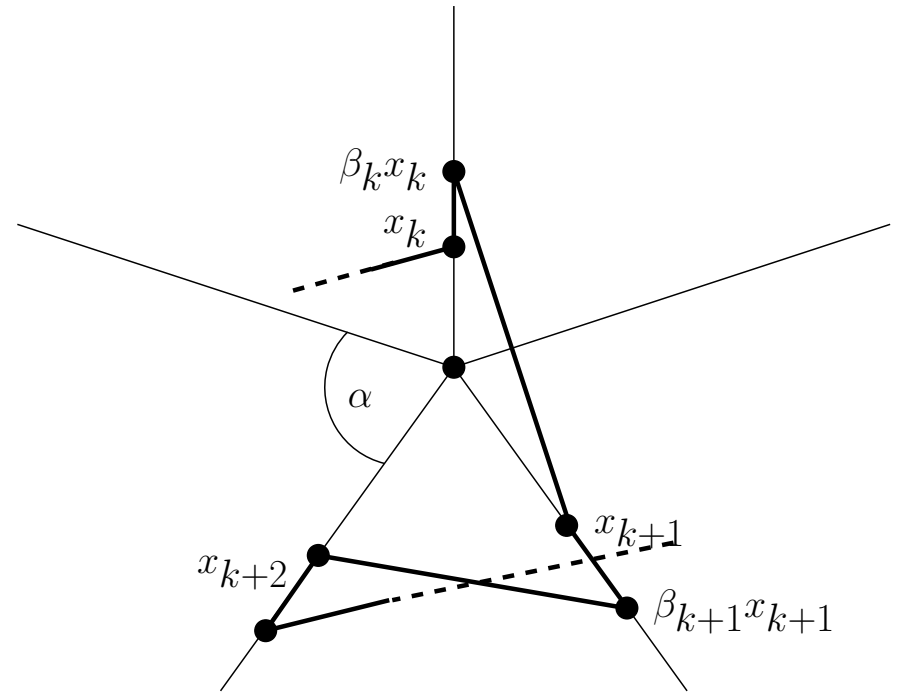
Lower bound: Special example

- General problem of lower bounds
- Special example: Searching for a special ray in the plane
- a special ray in the plane
- Cross R and detect s
- Special case: No move to s
- Alpern/Gal (Spiral search): Ratio $C = 17.289 \dots$
- Best ratio among monotone and periodic strategies
- "Complicated task": There is an optimal periodic/monotone strategy



Lower bound construction

- n known rays emanating from a ■
- Angle $\alpha = \frac{2\pi}{n}$ ■
- Find s on one of the rays ■
- Also non-periodic and non-monotone strategies ■
- Strategy S : Visits the rays in some order ■
- Hit x_k , leave $\beta_k x_k$ ($\beta_k \geq 1$) ■

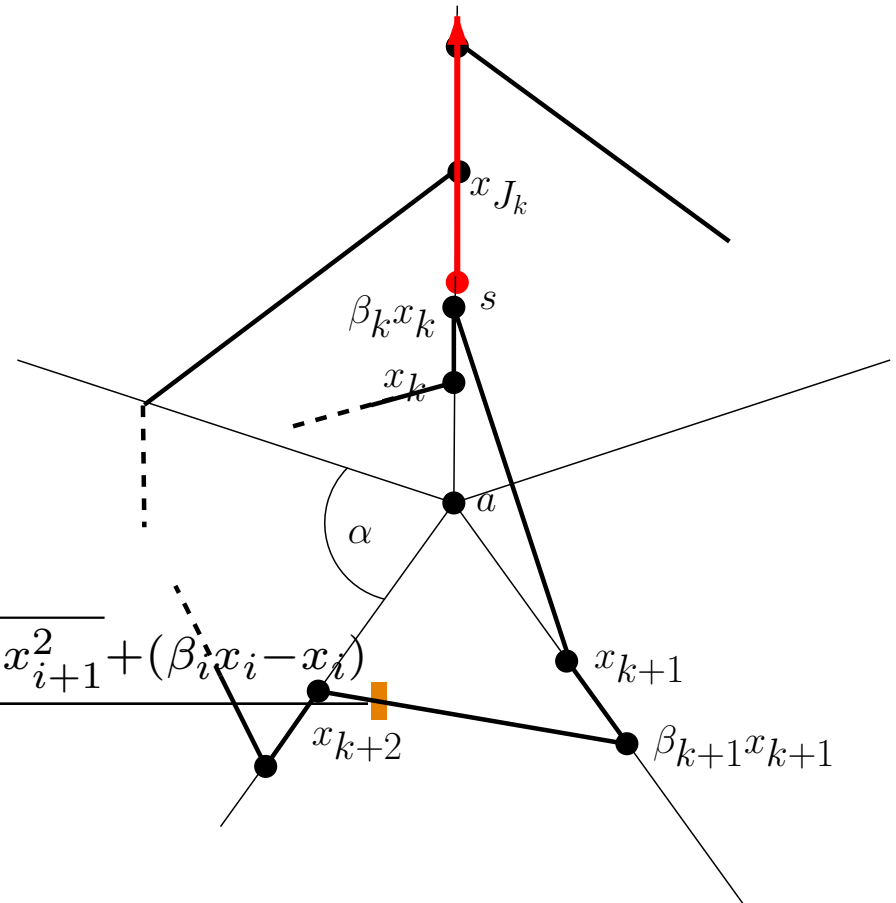


Lower bound construction

- Find s on a ray visited up to $\beta_k x_k$ at the last time, now at x_{J_k}
- Note: Any order is possible
- Worst-case, s close to $\beta_k x_k$
- Ratio: $C(S)$

$$\frac{\sum_{i=1}^{J_k-1} \sqrt{(\beta_i x_i)^2 - 2\beta_i x_i x_{i+1} \cos \gamma_{i,i+1} + x_{i+1}^2} + (\beta_i x_i - x_i)}{\beta_k x_k}$$

- Monotone/Periodic, Funktional??



- Ratio: $C(S)$ **Lower bound construction**

- $$\frac{\sum_{i=1}^{J_k-1} \sqrt{(\beta_i x_i)^2 - 2\beta_i x_i x_{i+1} \cos \gamma_{i,i+1} + x_{i+1}^2} + (\beta_i x_i - x_i) J_k x_{J_k}}{\beta_k x_k}$$

- Shortest distance to next ray:

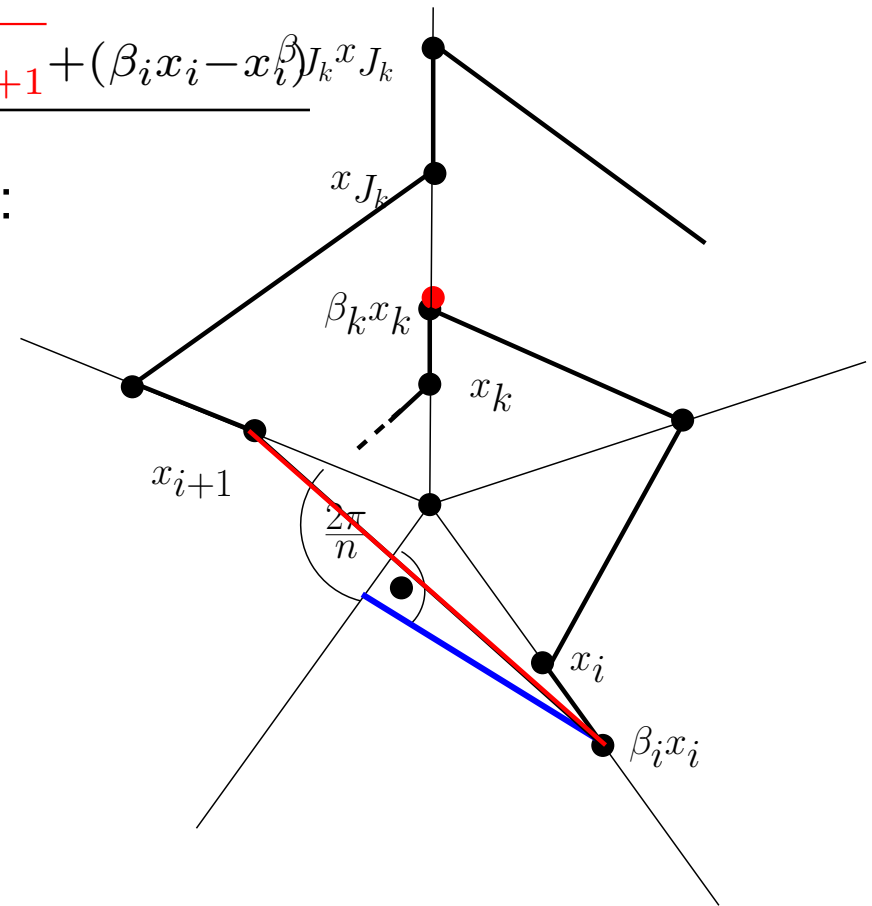
$$\beta_i x_i \sin \frac{2\pi}{n}$$

- Lower bound for

$$\sqrt{(\beta_i x_i)^2 - 2\beta_i x_i x_{i+1} \cos \gamma_{i,i+1} + x_{i+1}^2}$$

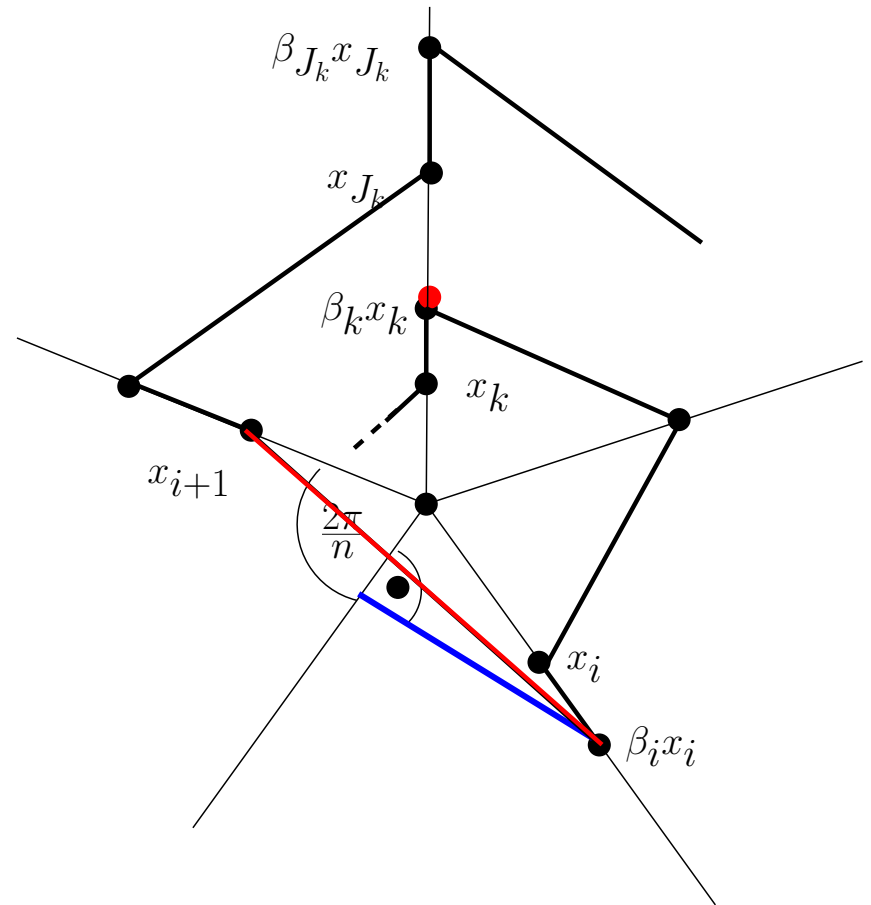
- Lower bound: $C(S) \geq$

$$\sin \frac{2\pi}{n} \frac{\sum_{i=1}^{J_k-1} \beta_i x_i}{\beta_k x_k}$$



Lower bound construction

- Lower bound $\frac{\sum_{i=1}^{J_k-1} f_i}{f_k}$ ■
- Equals functional of standard m-ray search ■
- Optimal strategy: monotone/periodic (Alpern/Gal) ■
- $f_i = \left(\frac{n}{n-1}\right)^i$
- ratio: $(n-1) \left(\frac{n}{n-1}\right)^n$ ■
- $C(S) \geq \sin \frac{2\pi}{n} (n-1) \left(\frac{n}{n-1}\right)^n$ ■

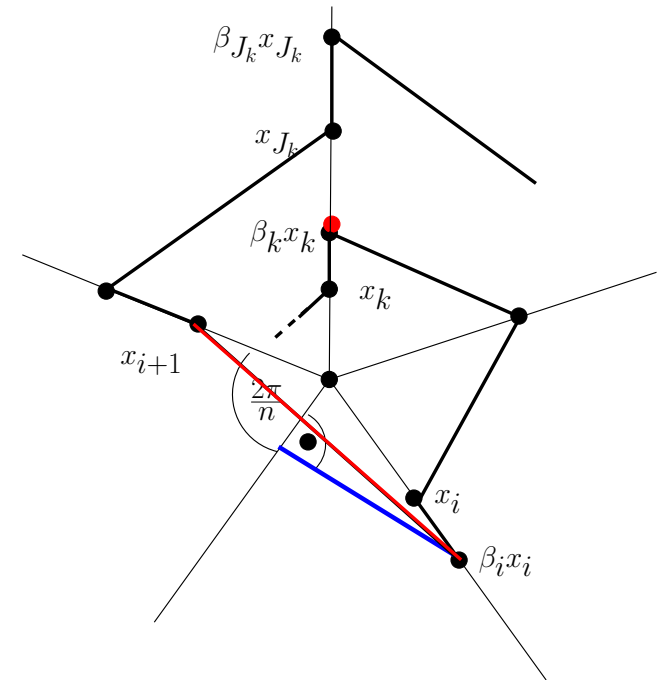


Lower bound construction

- $C(S) \geq \sin \frac{2\pi}{n} (n-1) \left(\frac{n}{n-1}\right)^n$ ■

- $$\lim_{n \rightarrow \infty} (n-1) \left(\frac{n}{n-1}\right)^n \sin \frac{2\pi}{n} =$$

- $$2\pi e = 17.079\dots$$



- Lower bound: **Theorem** ■

Summary

- The Window-Shopper-Problem
- Optimal strategy $C = 1.059 \dots$: **Theorem**
- Interesting design technique
- Rays in general
- Lower $C \geq 2\pi e = 17.079 \dots$ (**Theorem**) and upper bound $C = 22.51 \dots$ (**Theorem**)
- Lower bound construction
- Also a lower bound for special case with $C = 17.289 \dots$