

Case IV $a \in B$, let $B' := B \setminus \{a\} \subseteq A$

A sh. by $\tilde{\tau}_2 \Rightarrow \exists F \in \tilde{\tau}_2$ with $F \cap A = B'$

$\Rightarrow F, F \cup \{a\} \in \tilde{\tau}$

(Def. $\tilde{\tau}_2$) and

$$B = B' \cup \{a\} = (F \cup \{a\}) \cap (A \cup \{a\})$$

$\Rightarrow A \cup \{a\}$ sh. by $\tilde{\tau}$ and claim holds

$$\Rightarrow |\tilde{\tau}_2| \leq \binom{m-1}{0} + \dots + \binom{m-1}{d-1}$$

by ind. case d

Altogether:

$$|\tilde{\tau}| = |\tilde{\tau}_1| + |\tilde{\tau}_2|$$

$$\leq \binom{m-1}{0} + \binom{m-1}{1} + \dots + \binom{m-1}{d} +$$

$$+ \binom{m-1}{0} + \dots + \binom{m-1}{d-1}$$

$$= 1 + \binom{m}{1} + \dots + \binom{m}{d}$$

$$= \binom{m}{0} + \binom{m}{1} + \dots + \binom{m}{d}$$

from $\binom{m-1}{i} + \binom{m-1}{i-1} = \frac{(m-1)!}{i!(m-1-i)!} + \frac{(m-1)!}{(i-1)!(m-i)!}$

$$= \frac{(m-1)!(m-i) + (m-1)!i}{i!(m-i)!} = \frac{m!}{i!(m-i)!} = \binom{m}{i}$$

\Rightarrow Part (i) \checkmark

□

(ii) $\dim_{vc}(\mathbb{T}) \leq d$

Show: $\forall Y \subseteq X \quad |Y|=n$

$\mathbb{T}|_Y$ contains at most $\binom{n}{0} + \dots + \binom{n}{d}$ sets

Apply (i) on Y and $\mathbb{T}|_Y$ use $\dim_{vc}(\mathbb{T}|_Y) \leq \dim(\mathbb{T}) \leq d$ \square

(iii) $\binom{m}{0} + \binom{m}{1} + \dots + \binom{m}{d}$ has no closed form, due to knuth

Let $\frac{1}{2} \geq a > 0$ (f.x. a later)

$$1 = (a + (1-a))^m = \sum_{i=0}^m \binom{m}{i} a^i (1-a)^{m-i}$$

$$\geq (1-a)^m \sum_{i=1}^d \left(\frac{a}{1-a}\right)^i \binom{m}{i}$$

$$\geq (1-a)^m \sum_{i=1}^d \binom{m}{i} \left(\frac{a}{1-a}\right)^d$$

$$\left(\frac{a}{1-a}\right)^i \geq \left(\frac{a}{1-a}\right)^d$$

$$\Rightarrow \sum_{i=0}^d \binom{m}{i} \leq \frac{(1-a)^d}{(1-a)^m a^d} = S(a)$$

$$\left[\begin{array}{l} a \leq \frac{1}{2} \Leftrightarrow \\ \frac{a}{1-a} \leq 1 \end{array} \right]$$

$S(a)$ has minimum at $a = \frac{d}{m} \left(\leq \frac{1}{2}\right)$

$$S'(a) = \frac{(1-a)^d}{(1-a)^m a^d} \left(\ln(1-a) - \ln(a) \right)$$

< 0 non. increas. $a < \frac{1}{2}$

$$\sum_{i=0}^d \binom{m}{i} \leq \left(1 - \frac{d}{m}\right)^{d-m} \left(\frac{m}{d}\right)^d \quad (111)$$

$m \rightarrow \infty$

monotonically increasing

$$\lim_{m \rightarrow \infty} \left(1 - \frac{d}{m}\right)^{-m} = e^d$$

$$\begin{aligned} \ln(\dots) &= (d-m) \ln\left(\frac{m-d}{m}\right) \\ &= (d-m) (\ln(m-d) - \ln m) \end{aligned}$$

$$\lim_{m \rightarrow \infty} \left(1 - \frac{d}{m}\right)^d = 1$$

$\ln(\dots)' > 0$
for $d \leq m$

$$\begin{aligned} &\ln(m) - \ln(m-d) \\ &+ (d-m) \left(\frac{1}{m-d} - \frac{1}{m}\right) \end{aligned}$$

$$\Rightarrow \frac{d}{m} \leq \frac{1}{2} \quad \sum_{i=0}^d \binom{m}{i} \leq e^d \left(\frac{m}{d}\right)^d$$

Now $d \geq \frac{m}{2}$

$$f(d) = \left(\frac{e^m}{d}\right)^d$$

monotonically increasing w/d

$$\ln(f(d)) = d (\ln e^m - \ln d)$$

$$\ln(f(d))' > 0 \text{ for } d \leq m$$

$$\left(\ln m - \ln d \right) > 0$$

$\Rightarrow d \in \left[\frac{m}{2}, m\right]$:

$$\sum_{i=0}^d \binom{m}{i} \leq \sum_{i=0}^m \binom{m}{i} = 2^m < 2^{\frac{m}{2}} \cdot e^{\frac{m}{2}} = (2e)^{\frac{m}{2}} = f\left(\frac{m}{2}\right) \leq f(d) \quad \square$$

$2 < e$

Proof of Theorem 44 (ϵ -net Theorem) $\dim_{VC}(\mathcal{F}) = d < \infty, \epsilon > 0$
Existence of ϵ -net of size $C \cdot d \cdot \ln \frac{1}{\epsilon}$

Proof works as follows:

Choose set N of size $s = C \cdot d \cdot \ln \frac{1}{\epsilon}$ (also fix C)

M -randomly from X s independent random draws

$\Rightarrow M$ -probability that N is a $\frac{1}{\epsilon}$ -net is greater than 0. (So there has to be an ϵ -net. Magic)

Proof \rightarrow Existence, but also a bound for C

Preliminaries:

- Proof gives randomized algorithm for finding $\frac{1}{\epsilon}$ -net

- M -random sampling

Probability that next chosen point is $\in Y = M(Y)$

- s independent random draws, "day back"

Technical lemma distribution Series of size s not a set \emptyset

Lemma 46 Let $X = X_1 + X_2 + \dots + X_n$ where X_i

are independent random variables $\text{Prob}(X_i = 1) = p \quad \forall i$
 $\text{Prob}(X_i = 0) = 1 - p$

Then $\text{Prob}(X \geq \frac{1}{2}np) \geq \frac{1}{2}$ for $np \geq 8$

Deviation of value from exp. value/
Value of random variable
in a region around the expected value.

Proof: By Chebyshev's Inequality

$$P(|X - E(X)| \geq \alpha) \leq \frac{V(X)}{\alpha^2}$$

$$V(X) = E((X - E(X))^2)$$

$$E(X) = n \cdot p$$

$$V(X) = \sum_{i=1}^n V(x_i)$$

$$\stackrel{\text{indep.}}{=} \sum_{i=1}^n E((x_i - E(x_i))^2) \leq np$$

$$\begin{aligned}
 E((x_i - p)^2) &= E(x_i^2 - 2px_i + p^2) \\
 &= E(x_i^2) - 2pE(x_i) + p^2 \\
 &= p - p^2 \leq p
 \end{aligned}$$

$$E(x_i^2) = (1-p)0^2 + p \cdot 1^2 = p$$

$$E(x_i) = (1-p)0 + p \cdot 1 = p$$

complement

$$P(x < \frac{np}{2}) \leq P(|x - np| \geq \frac{1}{2} np)$$

$$= P(|x - E(X)| \geq \frac{np}{2}) \leq \frac{np}{(\frac{np}{2})^2} = \frac{4}{np}$$

$$P(x \geq \frac{np}{2}) \geq 1 - \frac{4}{np} \geq \frac{1}{2} \text{ for } np \geq 8. \quad \square$$

Now proof continued:

w.p.o.s. $M(F) \geq \frac{1}{r} \forall F \in \mathcal{F}$ (otherwise leave them out)

Two independent random draws N, M of size s ("lay back", sets)

Events:

E_0 : N is not an $\frac{1}{r}$ -net, i.e. $\exists S \in \mathcal{F} : S \cap N = \emptyset$
(one set is missed)

E_1 : $\exists S \in \mathcal{F} : S \cap N = \emptyset$ and $|S \cap M| \geq \frac{s}{2r}$ $\frac{s}{2r} := \frac{s}{2r}$

(Bound $\text{Prob}(E_0)$ from above by $\text{Prob}(E_1)$, $\text{Prob}(E_0) \leq 1$)

Claim I. $\text{Prob}(E_1) \leq \text{Prob}(E_0)$ clear E_1 requires more details \checkmark

Claim II. $\text{Prob}(E_0) \leq 2 \text{Prob}(E_1)$ ($E_0 \Rightarrow E_1$)

Proof (Claim II): Conditional probability

$\text{Prob}(E_1 | N)$ for fixed N , M is random

if N is $\frac{1}{r}$ -net $\Rightarrow \neg E_0 \wedge \neg E_1 \Rightarrow \text{Prob}(E_1 | N) = 0$
 $\text{Prob}(E_0 | N) = 0$

if N is not an $\frac{1}{r}$ -net

$\Rightarrow \exists S_N \in \mathcal{F}$ with $N \cap S_N = \emptyset$

$\Rightarrow P(E_0 | N) = 1$

$\Rightarrow P(E_1 | N) \geq P(|S_N \cap M| \geq \frac{s}{2r}) \geq \dots$

\uparrow
(some other set S' might also fulfill $S' \cap N = \emptyset$ and $|S' \cap M| \geq \frac{s}{2r}$)

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$$r = \frac{S}{2r} \quad M(S_N) \geq \frac{1}{r}$$

$$\Rightarrow r = \frac{S}{2r} \leq \frac{SM(S_N)}{2}$$

$$\dots \geq \text{Prob}(|S_N \cap M| \geq \frac{SM(S_N)}{2}) \geq \frac{1}{2}$$

$|S_N \cap M|$ random variable

$$\sum_{i=1}^S X_i \text{ drawing of } M$$

↑
independent ("lay back" is important here)

$$X_i = 1 \Leftrightarrow x \in S_N \text{ draw } x$$

Probability: $x \in S_N$ is $M(S_N)$

S-draws \Rightarrow Lemma 46 applicable
($SM(S_N) \geq 8, \frac{S}{r} \geq 8$)

Algorithm: $\forall N \quad \text{Prob}(E_0 | N) \leq 2 \text{Prob}(E_1 | N)$

for all fixed N .

$$\Rightarrow \text{Prob}(E_0) \leq 2 \text{Prob}(E_1)$$

$$\Gamma \quad P(A) = \sum_i P(B_i) \cdot P(A|B_i)$$

$$\Omega = \dot{\cup} B_i$$

total probability from partial probabilities \forall

Now different bound for $\text{Prob}(E_1)$ (reps $\text{Prob}(E_0)$ scale)

Therefore different drawing of N and M !

- Choose a series A of $2s$ elements of X ("lay back", prob. μ)
- Choose (randomly) s positions for N , the remaining ones build M

N, M are chosen with the same probability as before!

Let A be fixed. Consider fix $S \in \mathcal{F}$ (N, M randomly)

Let $P_S := P((N \cap S = \emptyset \text{ and } |M \cap S| \geq \alpha) | A)$

Claim III. $P_S \leq t \frac{-cd}{4} = \text{Prob}(E_1 | A)$
small!

Proof:

Case 1 $|A \cap S| < \alpha \Rightarrow |M \cap S| < \alpha \Rightarrow P_S = 0 \quad \checkmark$

Case 2 $|A \cap S| \geq \alpha$

$\Rightarrow P_S \leq P(N \cap S = \emptyset | A)$

only one part

$= P \left(\begin{array}{l} \text{The } s \text{ positions for } N \text{ out of } A \\ \text{avoid } (\geq \alpha) \text{ positions of } \\ S\text{-elements in } A \end{array} \right)$

("positions!")

possibilities for s different positions if x fixed and avoided

$$\leq \frac{\binom{2s-x}{s}}{\binom{2s}{s}} = \frac{(2s-x)! s! x!}{(s-x)! s! (2s)!} \quad \uparrow \frac{\text{pos. to avoid}}{\text{all}} \downarrow$$

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possibilities for s different pos. for fixed A / # random samples of s out of $2s$ positions

$$= \frac{(s-x_1)(s-x_2) \dots (2s-x)}{(s+1) \cdot (s+2) \dots (2s)}$$

$$\leq \left(\frac{2s-x}{2s}\right)^s = \left(1 - \frac{x}{2s}\right)^s = \left(1 - \frac{x}{2s}\right)^s \leq e^{-\frac{x}{2}}$$

$s \rightarrow \infty$

$$= e^{-\frac{cd \ln r}{4}} = e^{-\frac{\epsilon d}{4}}$$

$$\left(x = \frac{s}{2r}, s = cd \ln r\right)$$

$\dim_{\text{vc}}(\mathcal{F}) \leq d$ was not used up to now!

Sh. Func. Lemma

$$|A| \leq 2s \Rightarrow \mathcal{F}|_A \text{ has at most } \left(\frac{2es}{d}\right)^d$$

different elements.

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Let A be fixed but $S \in \tilde{T}$ variable

event $E(s) := (S \cap N = \emptyset \text{ and } |S \cap M| \geq r)$

depends on $A \cap S$

$$P[E_1 | A] \leq \sum_{\substack{A \cap S \\ \text{pairwise} \\ \text{different}}} \overbrace{P[E(s)]}^{P_S}$$

pairwise different

↑
" $\exists S \in \tilde{T} : E(s)$ "

$$\leq |\tilde{T}|_A \cdot \max_S P_S$$

$$\leq \left(\frac{2eS}{d} \right)^d \cdot r^{-\frac{Cd}{4}}$$

$$= \left(2eCr \ln r \cdot r^{-\frac{C}{4}} \right)^d < \frac{1}{2}$$

for $d, r \geq 2$ and sufficient large C .

Algorithm: $\text{Prob}(E_0) \leq 2 \text{Prob}(E_1) \leq 2 \sum_A \text{Prob}(E_1 | A) \text{Prob}(A) < \frac{1}{2} \underbrace{\sum = A}_{< 1}$

Probability that a randomly chosen N is not an $\frac{1}{r}$ net is smaller than 1 \Rightarrow Probability that a randomly chosen net is an $\frac{1}{r}$ net is bigger than 0 \Rightarrow There exists such a net!

Conclusions:

Art gallery: $\dim_{\mathbb{R}}(\widehat{\Gamma}) \leq 14$

$$\widehat{\Gamma} = \{ \int_{\mathbb{R}} \text{vis}_p(P) \mid p \in P \}$$

$$\text{vis}_p(P) \geq \frac{1}{10} \text{vol}(P) \quad \Gamma = 10$$

$$\Rightarrow d = 14 \quad \Gamma = 10$$

$$C \text{ so that } \left(2e C \cdot 10 \ln 10 \cdot 10^{-\frac{C}{4}} \right)^{14} < \frac{1}{2}$$

$$C = 13.08$$

$$S = \frac{C}{d} \cdot \Gamma \ln r = 13.08 \cdot 14 \cdot 10 \ln 10 = 4218 \text{ guards!}$$

(Art Gallery) Theorem 47 There is a constant $D > 0$

so that. For any area P bounded by a simple closed Jordan curve and any point in P sees at least $\frac{1}{r}$ of the volume of P the number of

guards required is smaller than $D \cdot r \ln r$!